

Motivation Questions

- How can we design audits when agents adversarially misreport their private type?
- Examples include
 - IRS:** Verifying tax deductions and self-reported income.
 - Insurance:** Checking claims on property or medical damages.
 - Market Surveillance:** Detecting fraudulent trading activities.

Summary

- Optimal Non-Adaptive Audit:** Algorithm finds an ε -optimal audit policy in $O(m^2)$ time.
- Social Welfare & Policy Levers:** An $O(m^2)$ algorithm for social welfare; moreover, higher penalties or lower audit costs weakly improve both the principal's utility and welfare.
- Learning without a prior:** We design a no-regret online method (EXP3-based) with Regret $O(n\sqrt{T}m^2 \log m)$
- Adaptive adds no value:** Best adaptive = best non-adaptive; the same $O(m^2)$ search.

Model

Agents

- n agents has a private type $i \sim q \in \Delta_m$. Each report type $k \in [m]$.
- Misreporting strategy** $Q \in [0,1]^{m \times m}$, where $Q_{i,k} = \Pr[i \rightarrow k]$
- The principal's audit policy $p \in [0,1]^m$
- Agents' utility

$$U_{i,k}(p) = \text{pay}(k) - p_k \text{pen}(i, k)$$

For instance, $\text{pen}(i, k) = 1[i \neq k](\text{pay}(k) + c)$ (affine penalty)

Principal's Strategy:

- Non-adaptive: commit a fix audit vector p ,
- Adaptive: $p = \pi(\hat{q})$ as a function of observed report distribution \hat{q}

Principal's utility ($n=1$)

- Revenue: $V(p, Q) = \sum_{i,k} q_i Q_{i,k} (\text{val}(i, k) - \text{pay}(k) + p_k \text{pen}(i, k))$
- Audit cost: $C(p, Q) = \sum_{i,k} q_i Q_{i,k} p_k$
- Costly setting:

$$V_\lambda(p, Q) = V(p, Q) - \lambda C(p, Q) \text{ with cost per audit } \lambda$$

- Budget setting:

$$V_B(p, Q) = \begin{cases} V(p, Q) & \text{if } C(p, Q) \leq B \\ -\infty & \text{otherwise} \end{cases}$$

- Agents best-respond to p (or π); the principal maximizes utility under the worst equilibrium (pessimistic/Wardrop).

$$\max_p \min_{Q \in Eqi(p)} V(p, Q)$$

Bayes-Nash Equilibrium (BNE)

- Q is (non-atomic) BNE under p if, for all types i and reports $k, l \in [m]$ with $Q_{i,k} > 0$:

$$U_{i,k}(p) \geq U_{i,l}(p)$$

We set $Eqi(p)$ as the set of BNE.

Non-adaptive audit

Theorem 1

For any small enough $\varepsilon > 0$ and any non-adaptive game with cost λ and parameters $(n, m, q, \text{val}, \text{pay}, \text{pen})$, there is an algorithm that computes a $2n\varepsilon$ -optimal audit policy in $O(m^2)$ time.

Welfare objective: The same audit-search methods.

Monotonicity: Higher penalties or lower audit costs improve utility and welfare.

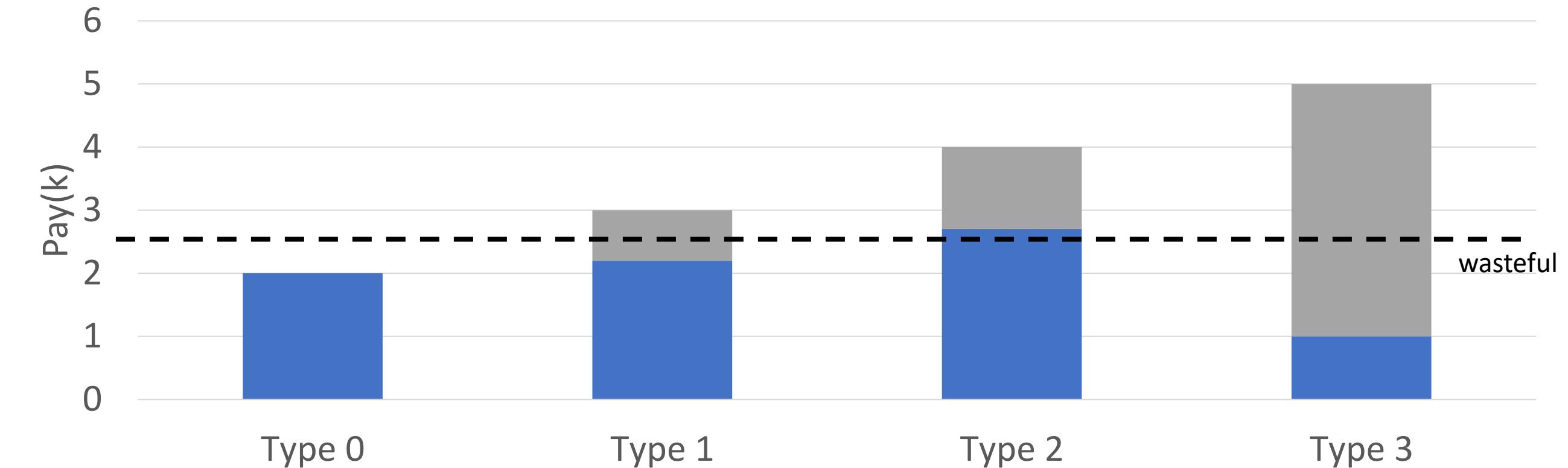
Relaxed priors: Even without exact knowledge of q , a no-regret online learner achieves near-optimal audits.

Equilibrium Characterization

- Idea: agents' best response is a threshold strategy

$$\hat{U}_k = \text{pay}(k) - p_k \text{pen}(k)$$

misreport truthful



- Choose the audit vector that **equals** \hat{U}_k for all k

$$\rho_k(u) = \frac{\text{pay}(k) - u}{\text{pen}(k)}$$

- The principal chooses from a family of audit policy, $p = \text{equa}(u, A, \varepsilon)$ to make agents' highest misreport utility is u :

$$p_k = \begin{cases} 0 & k < \iota \\ \rho_k(u) & k \in A \\ \rho_k(u - \varepsilon) & \text{otherwise} \end{cases}$$

Adaptive setting

Assumption on the penalty

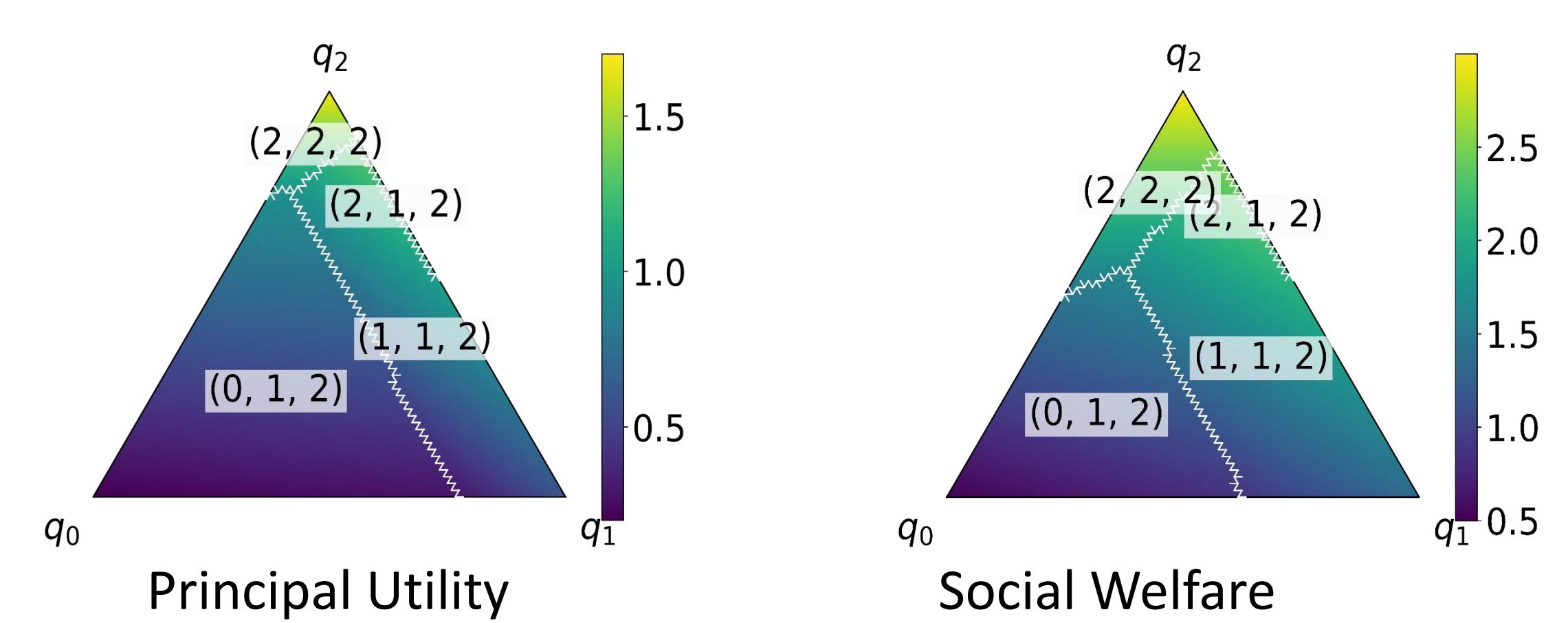
- Type independent: $\text{pen}(i, k) = 1[i \neq k]\text{pen}(k)$
- Insensitive: $\frac{\text{pay}(l)}{\text{pay}(k)} \geq \frac{\text{pen}(l)}{\text{pen}(k)}$

Theorem 2 (informal). Under the type independent and insensitive condition, adaptive audits do not outperform non-adaptive ones:

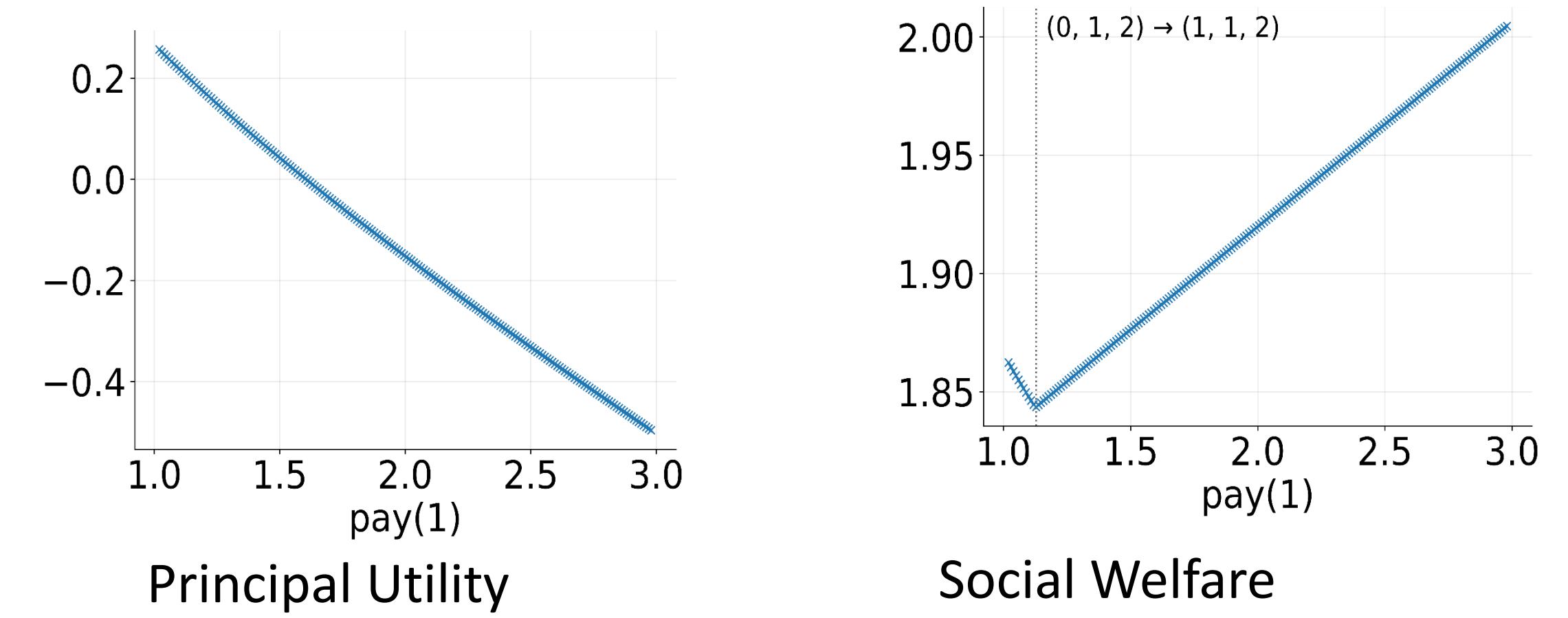
$$\max_{\pi} \min_{Q \in Eqi(\pi)} V_\lambda(\pi, Q) = \max_p \min_{Q \in Eqi(p)} V_\lambda(p, Q).$$

Simulations

Effect of prior



Effect of payment



Future work

- Generalize to finite agents, noisy or partial verification, and richer penalty structures.
- Design problem of payment and penalty function.
- Connect to **security games** and **tolling games**