

From Crowds to Codes: Minimizing Review Burden in Conference Review Protocols

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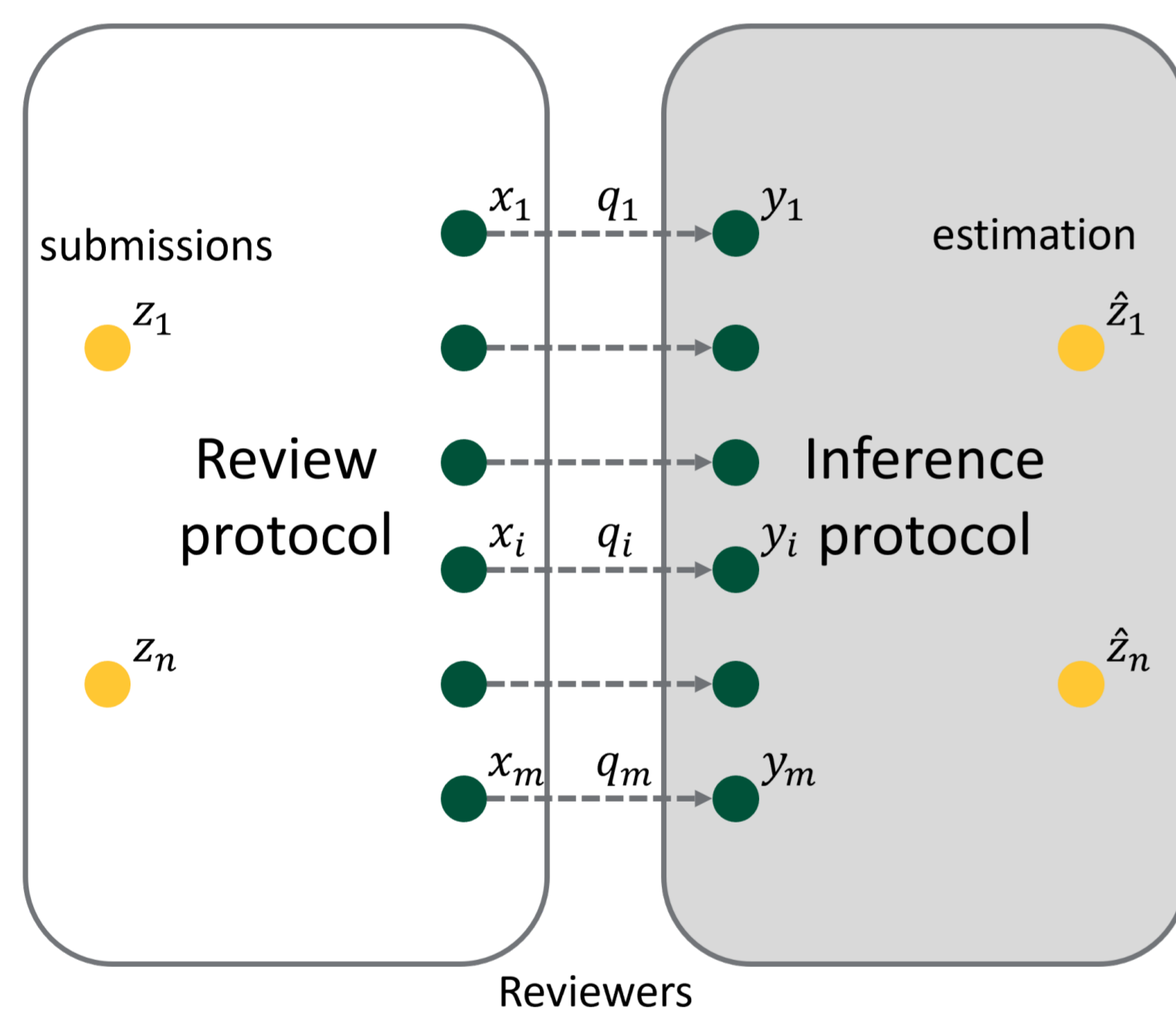
Key Contributions

How can we minimize the *review load ratio*—the number of review tasks per paper—while ensuring error probability at most ϵ ?

- Isolated Protocols: Optimal load ratio is $\Theta(\ln n/\epsilon)$ → grows with number of papers n and desired accuracy.
- Joint Protocols: Optimal load ratio is a constant (depends only on reviewer noise) → Joint protocols, e.g., two-phase review, can reduce the review burden.
- Unknown Noise Levels: We analyze how to design efficient protocols without knowing reviewer noise levels.

Conference Review Problem

A conference review problem consists of the number of paper n and the distribution of noise level \mathcal{D}_q .



Review and Inference Protocols

Review Protocols

We can classify the review protocols based on two key attributes: *isolated* vs *joint* and *parallel* vs *adaptive*.

- A protocol is isolated if each review task pertains to a single paper.
- A protocol is non-adaptive (or parallel) if review tasks are independent of each other and the reports are simultaneously collected.

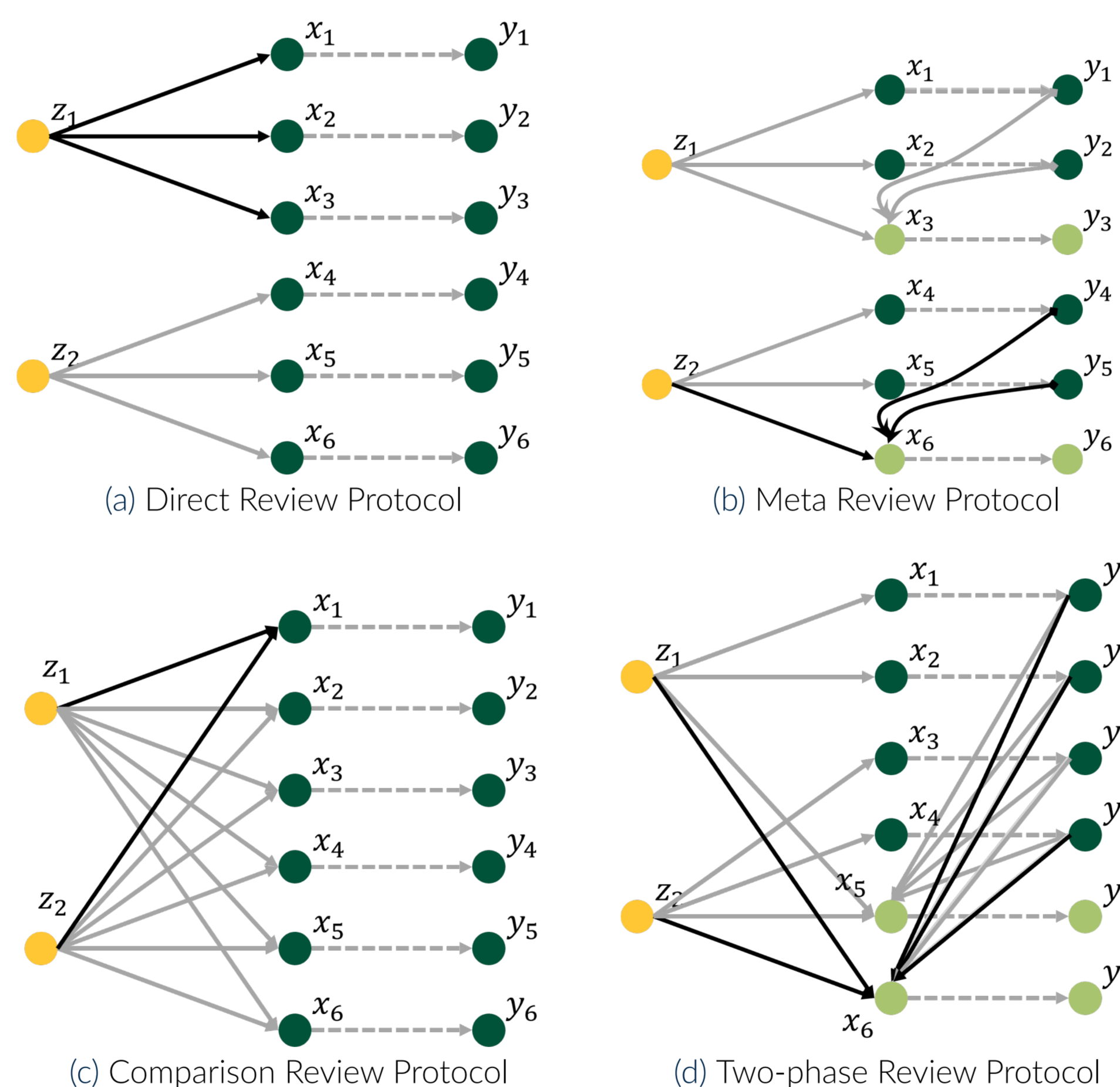


Figure 1. Examples of review protocols and tasks

Examples:

- A task i is a direct review task of paper j , if the report is a noisy estimate of paper j 's quality $y_i \stackrel{q_i}{\leftarrow} f_i(z) = z_j$ and the true answer is $x_i = z_j$
- A task i is second phase review task can be

$$y_i \stackrel{q_i}{\leftarrow} f_i(z, \mathbf{y}_{<i}) = \begin{cases} z_1 & \text{if } y_1 + y_2 \geq y_3 + y_4 \\ z_2 & \text{otherwise.} \end{cases}$$

Inference Protocols

The conference cares about the *error probability*, $\Pr[\hat{z} \neq \mathbf{z}]$. Given a review protocol \mathbf{f} , the optimal inference rule minimizing this error is the maximum a posteriori (MAP) estimator.

A (λ, ϵ) -protocol has the load ratio is upper bounded by λ and error probability is upper bounded by ϵ .

Minimizing the Review Load Ratio

Isolated Protocols

Given (n, \mathcal{D}_q) and small enough $\epsilon > 0$, an optimal isolated protocol with error probability $P_e = \epsilon$ requires load ratio $\lambda = \Theta(\ln \frac{n}{\epsilon})$.

- The reports of direct Review Protocols Blackwell dominates any isolated protocols'.
- Compute the optimal load ratio of direct review protocols.

Joint Protocols

Let $h(q) := q \log_2 \frac{1}{q} + (1-q) \log_2 \frac{1}{1-q}$ and $\lambda^* := \left(1 - \sum_{q \in \text{supp}(\mathcal{D}_q)} \mathcal{D}_q(q) h(q)\right)^{-1}$.

For any $\delta, \epsilon > 0$, there is no $(\lambda^* - \delta, \epsilon)$ -adaptive protocol, but there exists a $(\lambda^* + \delta, \epsilon)$ -nonadaptive conference protocol, when the number of papers n is sufficiently large.

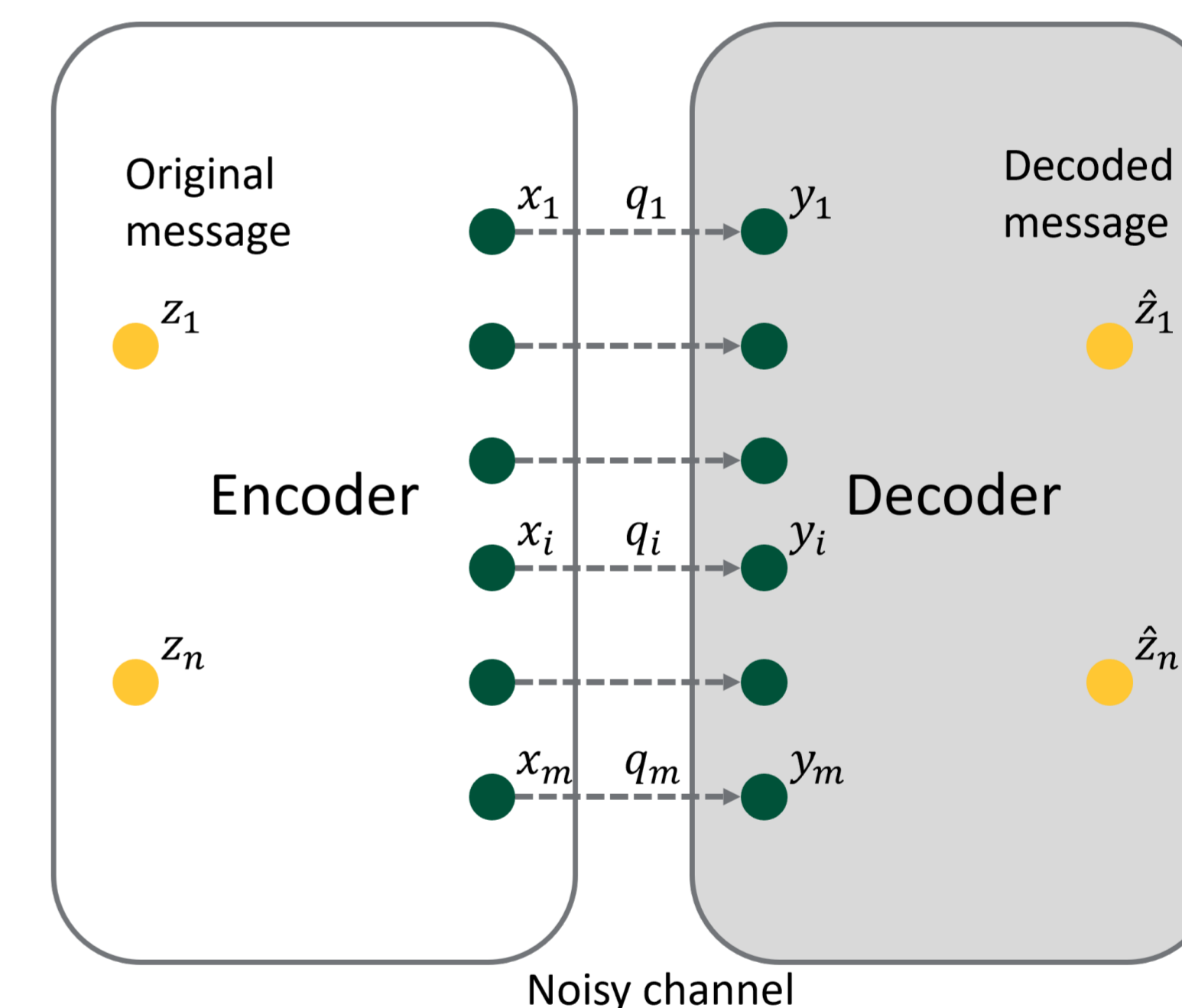


Figure 2. Connection to coding theory

- n submissions with quality $z_j \in \{-1, 1\}$ for all $j \in [n]$ uniformly distributed.
- m review tasks each has noise level $q_i \in [0, 1]$, binary question f_i with answer x_i , and a report y_i in $\{-1, 1\}$ so that

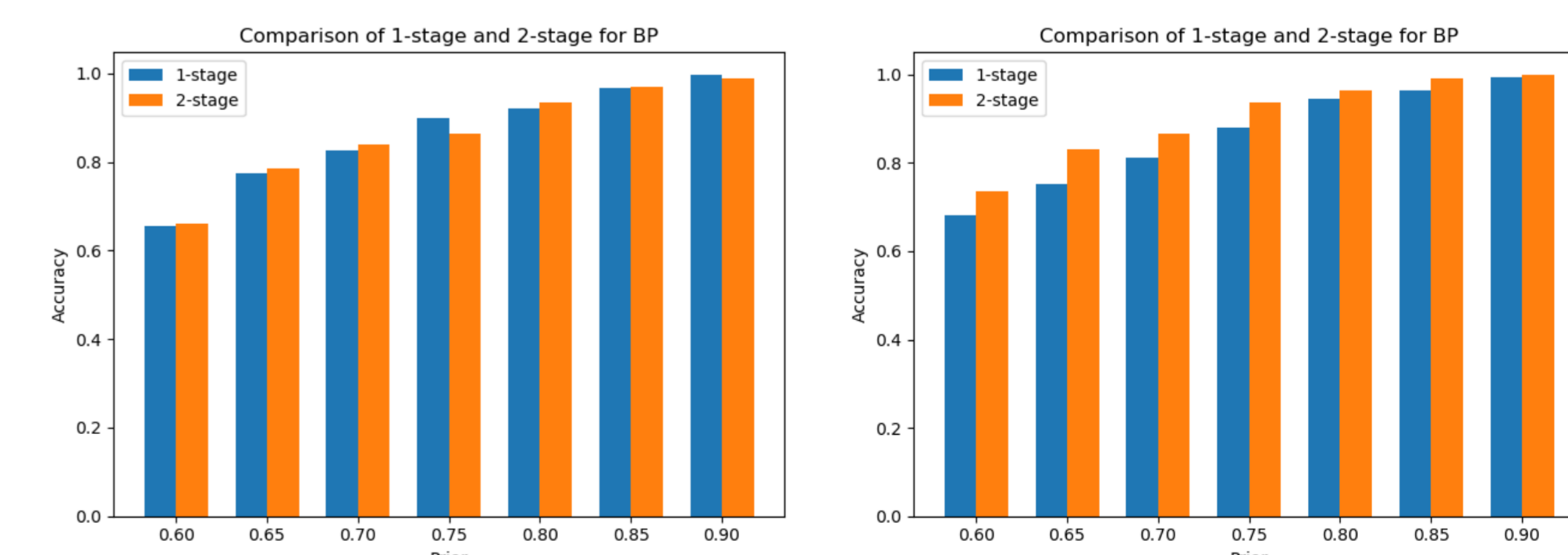
$$y_i = \begin{cases} x_i & \text{with probability } 1 - q_i \\ -x_i & \text{otherwise.} \end{cases}$$

- Histogram of noise level \mathcal{D}_q : quality of PC pool
- Review load ratio $\lambda := \frac{m}{n}$: number of review tasks per paper

Given (n, \mathcal{D}_q) , a conference protocol with load ratio λ consists of

- **Review protocol:** Design $m = \lambda n$ review tasks $\mathbf{f} = (f_1, \dots, f_m)$, and agents submit reports $\mathbf{y} = (y_1, \dots, y_m)$.
- **Inference protocol:** Infer the quality of n papers, $\hat{\mathbf{z}}$, using \mathbf{y} .

Simulations



(a) Select the bottom 50% most inconsistent papers for the second phase (b) Select the top 75% of papers based on aggregated scores for the second phase

Figure 3. Two-phase review with 1/4 review tasks in the second phase