

Research Challenges in Securing VoIP



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What Today's Talk is All About?

- Voice over IP (VoIP) basics
 - Signaling, voice stream, billing
- Security threats to VoIP
 - What are they?
 - How real are they?
 - Current VoIP security mechanisms
- Why securing VoIP is so challenging?
 - Open architecture, it's Internet
 - No prior established trust between caller & callee
 - Key management challenges
- Objectives of this tutorial
 - Bring the attention of the research community to the problems of VoIP security
 - Discuss the research challenges and open problems in securing VoIP
 - Seek your insight on securing VoIP



Outline

- Introduction
 - Motivation
- VoIP basics
 - Signaling, voice stream, billing
 - Security mechanisms
- Threats to VoIP security
 - Registration hijacking
 - DoS
- Exploits of VoIP security
 - Billing attack
 - Nuisance call
- Mitigations to VoIP exploits
 - IDS, cross-protocol correlation
- Research challenges and open problems
 - NAT traversal, key management

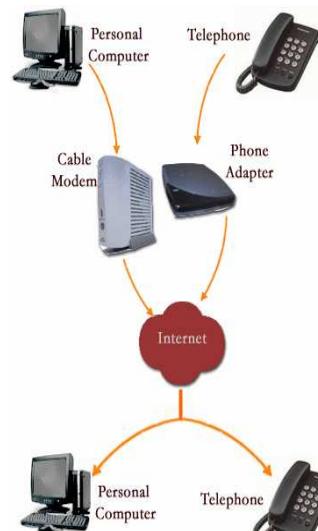


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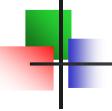
Proliferation of VoIP

- People are moving from POTS to VoIP service
 - VoIP is cheaper, more convenient and flexible, and it provides more features
- The number of residential VoIP subscribers worldwide is expected to grow from
 - current 38 million
 - to 267 million by 2012
- Types of VoIP
 - Managed
 - Vonage, AT&T, Verizon etc.
 - Unmanaged
 - Computer to computer VoIP (e.g., Skype)



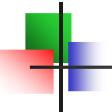
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Expectations on VoIP

- Confidentiality Security – No.1 concern
- E-911
- Lawful surveillance
- QoS



Security Requirements on VoIP

- **Authenticity**
 - When A dials B's number, the call will reach B.
 - The incoming call is really from who it claims to be – caller ID is authentic
- **Confidentiality, privacy and anonymity**
 - No one other than the caller(s) and callee(s) should
 - have access to the conversation content
 - Even know that Alice and Bob have talked over VoIP
- **Integrity**
 - The call signaling and content have not been tampered with
- **Temper resistant billing**
 - Service provider – no service stealing, toll fraud
 - Subscriber – no overcharge
- **Availability**
 - be resilient to denial-of-service (DoS) attack





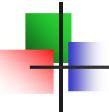
Key Functional Components of VoIP

- VoIP signaling
 - Responsible for establishing, managing, tearing down VoIP sessions
 - H.323 – 1996 ITU
 - MGCP – 1999 IETF RFC
 - **SIP (Session Initiation Protocol)** – 1999, 2002 IETF RFC
 - The dominant VoIP signaling protocol
- VoIP voice stream
 - RTP (Real-time Transport Protocol)
 - SRTP (Secure RTP)
- VoIP billing
 - Indispensable for VoIP service providers (e.g., Vonage, AT&T, Vonage)



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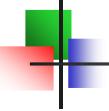
SIP Overview

- SIP is an IETF standard
 - RFC 2543 of 1999 (obsolete)
 - RFC 3261 of 2002
- SIP is
 - “an application-layer control (signaling) protocol for creating, modifying, and terminating sessions with one or more participants” (RFC 3261)
 - Text based, very similar to HTTP
- SIP sessions include
 - Internet telephone call
 - Multimedia conferences
 - Instant messaging
 - etc.



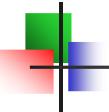
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SIP Functionalities

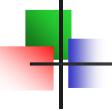
- Learn, determine the location, availability of remote communicating party
 - Registration
 - Call routing
 - Call redirection
- Establish, a session between end points
 - Call setup, transfer, termination
- Negotiate the media capability
 - Use SDP (Session Description Protocol) to specify the media parameters (e.g., IP address, port number, codec)



SIP Components

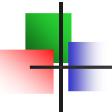
- User Agents – owned/used by subscribers
 - User agent client (UAC) – who initiates a call
 - User agent server (UAS) – who receives a call
- SIP Servers – maintained by service providers
 - Proxy server
 - relays the signaling messages (and potential voice streams) between the caller and callee
 - Location server
 - Keeps where a subscriber's UA is currently at (the IP address)
 - Registrar server
 - Accept registration from subscribers about their current locations
 - Keeps track subscribers' whereabouts at Location server
 - Redirect server
 - Provides information about next hop to a subscriber





SIP Messages

- Two types
 - Request – identified by a method name
 - Response – identified by a number similar to HTTP
- SIP request messages
 - REGISTER – tell where the UAC is currently at
 - INVITE – initiate a call to someone
 - BYE – terminate an established call
 - ACK – acknowledge the receipt of some message
 - CANCEL – quit from an ongoing call setup
 - OPTION – to query the capability of a server

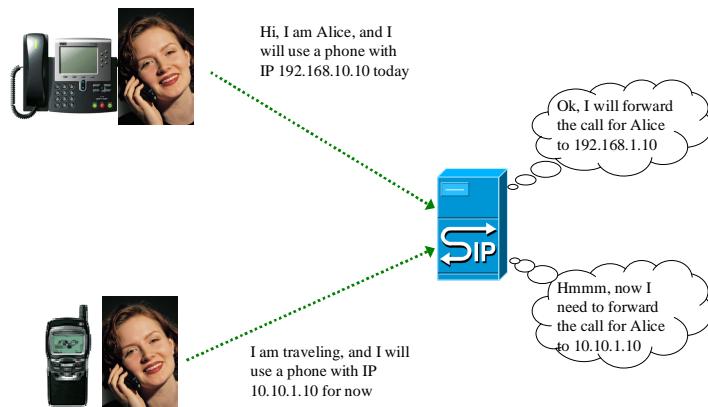


SIP Messages (cont'd)

- SIP response messages
 - 1xx Provisional – 100 Trying, 180 Ringing
 - 2xx Successful – 200 Ok
 - 3xx Redirection – 301 Moved Permanently, 302 Moved Temporarily
 - 4xx Failure – 404 Not Found, 410 Gone, 403 Forbidden
 - 5xx Server Failure – 503 Service Unavailable
 - 6xx Global Failure – 600 Busy Everywhere



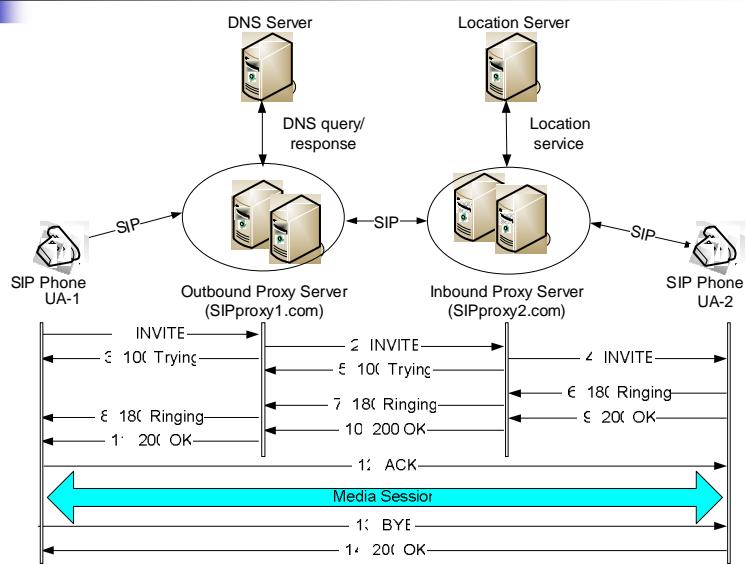
SIP Registration



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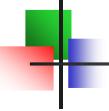
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SIP Message Flow of Normal Call Setup & Tear Down



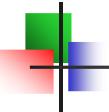
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VoIP Security

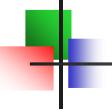
- VoIP signaling security
 - Protect the authenticity, integrity of the signaling message
- VoIP voice stream security
 - Protect the authenticity, integrity and confidentiality of the voice content (including any keys pressed)
- VoIP billing security
 - Prevent service theft, toll fraud, undercharge, overcharge



SIP Security Mechanisms

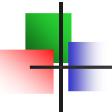
- SIP does NOT define its own security mechanism, it reuses existing security mechanisms for HTTP, SMTP whenever possible
- Two building blocks of SIP security mechanisms
 - Authentication
 - [HTTP digest authentication](#)
 - Encryption
 - [IPsec, TLS, S/MIME](#)
- SIP authentication and encryption can NOT be applied to the whole SIP messages from end-to-end
 - Intermediate SIP proxies need to modify and insert SIP message fields
 - Add via field
 - Change request URI due to call-redirection
- SIP security is hop-by-hop





HTTP Digest Authentication

- Provides
 - one-way authentication
 - Identify UAC to a UAS or SIP Proxy
 - Does NOT identify UAS or SIP proxy!
 - anti-replay protection
- Assumes the two parities involved share a secret password
 - Usually hard-coded in the UAC (SIP phone adaptor)
- Uses challenge/response
 - Response = $F(\text{nonce}, \text{username}, \text{password}, \text{realm}, \text{SIP-method}, \text{request-URI})$
- Must be supported by all SIP compliant UA and SIP servers

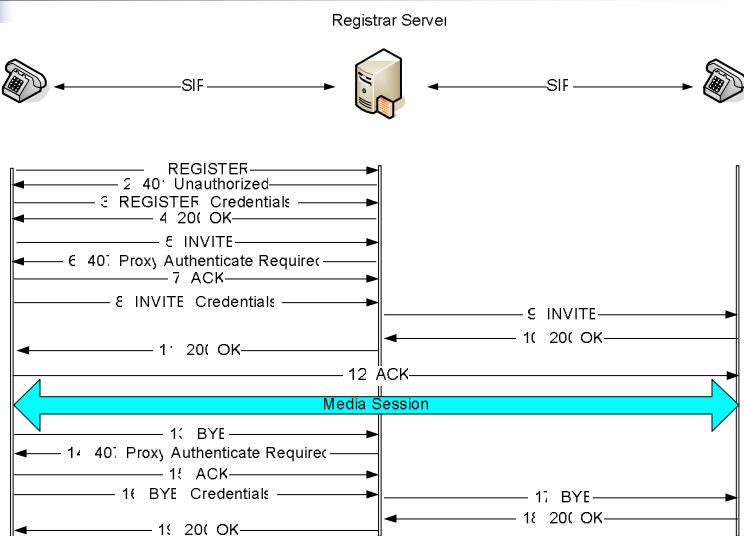


SIP Security Mechanisms (cont'd)

- TLS
 - Can provide authentication, integrity, confidentiality
 - But authenticate the server only
 - Applied hop-by-hop
 - All SIP compliant servers MUST support TLS
 - UAs are strongly recommended to support TLS
 - Each UA (SIP) has its own certificate?
- IPsec
 - Can provide authentication, integrity, confidentiality
 - No SIP component is required to support IPsec
 - Can be applied hop-by-hop
 - Key management issues
 - Setting up security association with every UA is too expensive
- S/MIME
 - Can provide some degree of end-to-end authentication, integrity or confidentiality for most SIP header fields
 - Excluding Request-URI, Via, Record-Route, Route, Max-Forwards, and Proxy-Authorization
 - Require PKI to be effective
 - No SIP component is required to support S/MIME



Message Flow of SIP Authentication



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Unauthenticated INVITE Message

```

Session Initiation Protocol
Request line: INVITE sip:*****3255@d.voncp.com:10000 SIP/2.0
Method: INVITE
Message Header
Via: SIP/2.0/UDP 192.168.0.108:10000;branch=z9hG4bK-85c98d52
From: <sip:*****3953 <sip:1*****3953@d.voncp.com:10000>;tag=6297ece2276a64060
To: <sip:*****3255@d.voncp.com:10000>
Remote-Party-ID: *****-3953 <sip:1*****3953@d.voncp.com:10000>;screen=yes;party=calling
Call-ID: a6b81312-da84f396@192.168.0.108
CSeq: 101 INVITE
Max-Forwards: 70
Contact: <*****-3953 <sip:1*****3953@192.168.0.108:10000>
Expires: 240
User-Agent: 0013101DCFB Linksys/RT31P2-3.1.6(LI)
Content-Length: 308
Allow: ACK, BYE, CANCEL, INFO, INVITE, NOTIFY, OPTIONS, REFER
Supported: x-sipura
Content-Type: application/sdp
Session Description Protocol
Session Description Protocol Version (v): 0
Owner/Creator, Session Id (o): - 2305 2305 IN IP4 192.168.0.108
.
.
.
Connection Information (c): IN IP4 192.168.0.108
.
.
.
Media Description, name and address (m): audio 10076 RTP/AVP 2 0 8
18 100 101

```



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407 Authentication Required Message

```
Session Initiation Protocol
Status line: SIP/2.0 407 Proxy Authentication Required
Status-Code: 407
Message Header
Via: SIP/2.0/UDP 192.168.0.108:10000;branch=z9hG4bK-85c98d52
From: ***-***-3953 <sip:1*****3953@d.voncp.com:10000>;
tag=6297ece2276a640600; natted=xx.***.102.135
To: <sip:*****3255@d.voncp.com:10000>
Call-ID: a6b81312-da84f396@192.168.0.108
CSeq: 101 INVITE
Proxy-Authenticate: Digest realm="69.59.227.87",
domain="sip:69.59.227.87", nonce="2036652154", algorithm=MD5
Max-Forwards: 15
Content-Length: 0
```

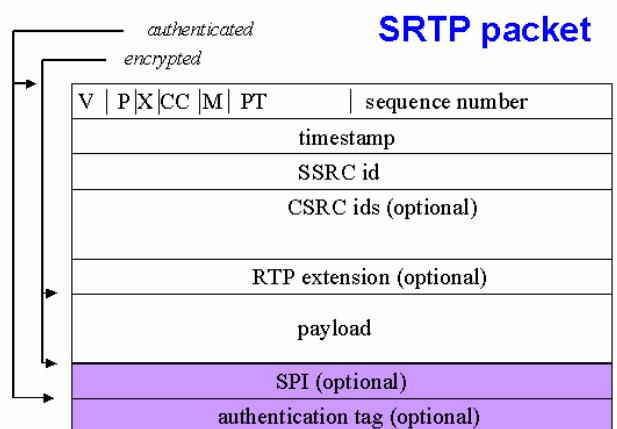


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<sip:1*****3953@d.voncp.com:10000>;tag=6297ece2276a640600
To: <sip:*****3255@d.voncp.com:10000>
Remote-Party-ID: ***-***-3953
<sip:1*****3953@d.voncp.com:10000>;screen=yes;party=calling
Call-ID: a6b81312-da84f396@192.168.0.108
CSeq: 102 INVITE
Max-Forwards: 70
Proxy-Authorization: Digest username="1*****3953",
realm="69.59.227.87", nonce="2036652154",
uri="sip:*****3255@d.voncp.com:10000", algorithm=MD5,
response="690680e4e138b38c1ba95271cc691b47"
Contact: ***-***-3953 <sip:1*****3953@192.168.0.108:10000>
Expires: 240
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18 100 101
```



Secure RTP



Elisabetta Carrara

SRTP

Image from IETF proceedings



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Threats to VoIP Security

- Service stealing
 - Steal minutes from VoIP service provider
 - Call at other subscriber's expense
- Service disruption against
 - the VoIP infrastructure
 - individual subscriber
 - terminate established calls
 - prevent certain calls from being established
- Call hijacking
 - Registration spoofing
 - Unauthorized call redirection
 - Taking over established VoIP call via re-INVITE
- Interception and modification
 - Conversation alteration – mix, change the RTP voice stream



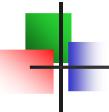
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Threats to VoIP Security (cont'd)

- Eavesdropping
 - Wiretap and traffic analysis
- VoIP fraud
 - Voice phishing (aka vishing)
- Annoyance
 - SPIT (Spam over Internet Telephony)
 - Nuisance call
- Attack against others
 - Divert VoIP traffic to flood someone
- VoIP based botnet
 - What if attacker compromises the softphone running in the laptop or handheld?



VoIP Service Stealing Case

- In June 2006, FBI arrested Pena and Moore for VoIP fraud
 - They
 - broke into networks of 15 VoIP service providers and
 - routed calls through them
 - sold the stolen (up to 10million) minutes at rate as low as 0.4 cents a minute to telecommunications providers
 - One victim VoIP service provider handled 500,000 calls in 3 weeks
 - <http://vonmag.com/editorial/web-exclusives/voip-fraud-hack>



SIP Phone Registration Spoofing

- Before a SIP UA (i.e. phone) can be used to make or receive a call, it must register itself so that SIP proxy server knows its current IP address.
- What if some attacker sends a spoofed REGISTER message to the registrar to trick the registrar into believing that the victim SIP phone is at an IP address chosen by the attacker
 - All the calls to the victim will reach the attacker's phone!
- What about SIP digest authentication?
 - The SIP registrar does not require authentication for registration
- Assuming the attacker does not know the secret password shared between the victim SIP phone and the registrar, can the attacker still spoof registration
- Yes!

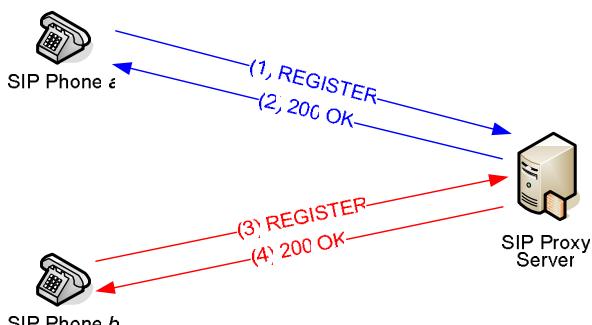


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SIP Phone Registration Spoofing

- The SIP digest does NOT cover the IP address of the SIP phone!
- The registrar actually uses the source IP address of the packet containing the REGISTER message
- The attacker could simply replay the legitimate REGISTER message from a different IP address

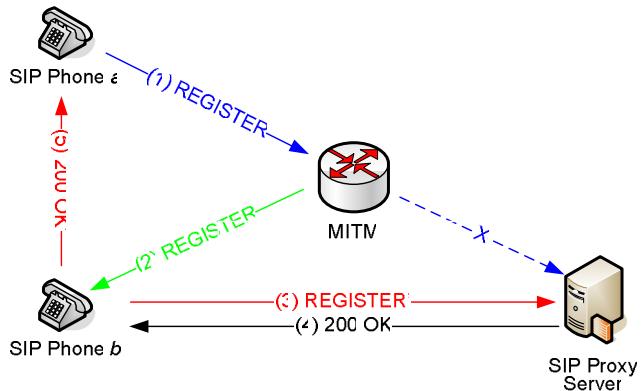


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SIP Phone Registration Hijacking

- With a MITM, attacker could completely hijack calls to the victim

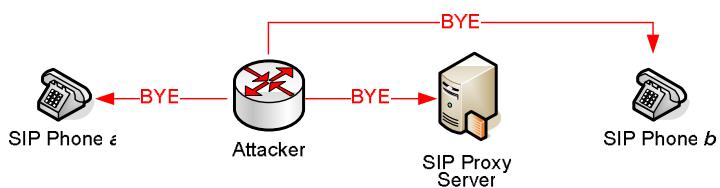


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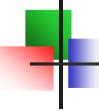
Fake BYE to Terminate Established SIP Call

- Attacker could terminate any established SIP call by sending fake BYE message to the SIP phone(s) and/or SIP proxy
- SIP proxy requires digest authentication, so it could detect fake BYE and ignore it.
- Existing SIP phones however will honor any BYE message with correct Call-ID.



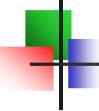
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VoIP CallerID Spoofing

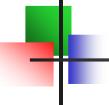
- CallerID of SIP call is NOT trustworthy, and it is easy to spoof
 - Just modify the From field of INVITE message
- There are companies that offer callerID spoofing service to the public
 - <http://www.telespoof.com/>
 - <http://www.spoofcard.com/>
 - <http://www.spoofcom.net/>
 - <https://www.itellas.com/>
 - <http://www.spooftel.com/>
- More information can be found at
<http://www.calleridspoofing.info/>



Billing of Managed VoIP

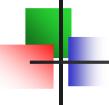
- Billing is fundamental to VoIP service providers (e.g. Vonage, AT&T)
 - Certain VoIP calls are charged on a per minute basis
 - International call
 - 900 call
 - Service providers rely on accounting and billing for charging their customers for the service they provided
 - ✖ Loss any revenues from any billable services they provide
 - ? Overcharge to customers
- Billing has direct impact on each individual VoIP subscriber
 - ✓ Charges for the VoIP services they have chosen and used
 - ✖ Charges on the calls they have NOT made
 - ✖ Overcharged on the calls they have made





Requirements of VoIP Billing

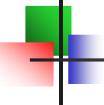
- Billing needs to be reliable
 - Resilient to billing fraud
 - Provides consistent view on what the service provider has provided and what the subscriber has received
- Billing needs to be trustworthy
 - It will determine
 - how much money the service provider will make
 - how much money the subscriber will pay
 - Inaccurate or corrupted bill will create disputes between the service provider and its customers



Billing and Signaling of VoIP

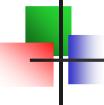
- Existing VoIP billing is based on VoIP signaling
 - Signaling determines
 - the caller and callee of the call
 - when the call starts and ends
 - where the call will be routed
- Any vulnerability in VoIP signaling could be a potential vulnerability of VoIP signaling
- VoIP signaling protocols
 - SIP – the dominant VoIP signaling protocol
 - MGCP
 - H323
- We will focus on billing of SIP-based VoIP





VoIP Billing Vulnerabilities

- Any vulnerability in VoIP signaling could be a potential vulnerability of VoIP signaling
 - Manipulation of SIP messages
 - Fake SIP messages
- The use public Internet for signaling and billing opens many doors for attacks
 - MITM at any router or gateway along the VoIP signaling path
- How vulnerable are those deployed commercial SIP-based VoIP services?
 - Vonage
 - AT&T
 - Verizon
 - Broadvoice



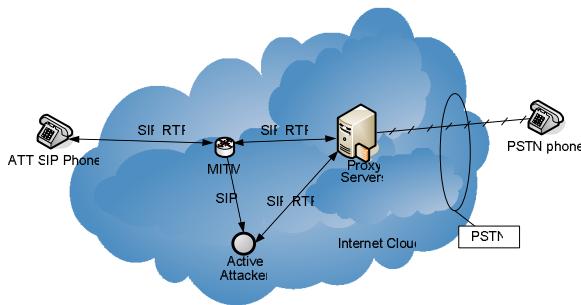
What We Have Done?

- Examined the billing vulnerabilities of SIP-based VoIP
- Identified a number of billing attacks on subscribers of deployed VoIP services
 - InviteReplay
 - FakeBusy
 - ByeDelay
 - ByeDrop
- Experimented with 2 leading VoIP services in US
 - Vonage – no. 1 in market share 53.9%
 - AT&T callvantage – no.2 in market share 5.5%



INVITE Replay Billing Attack

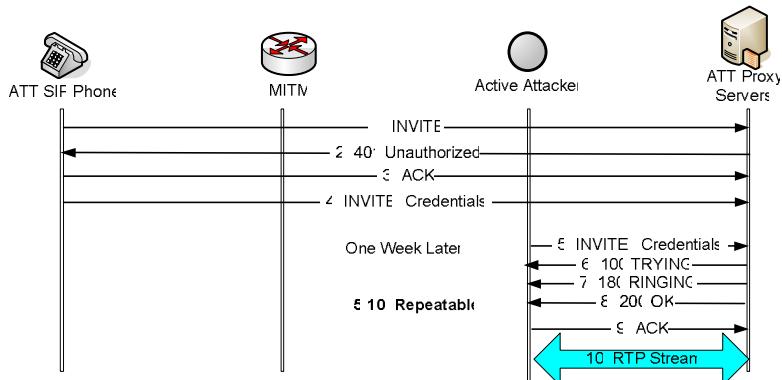
- Caller's SIP phone initiates a call by sending an INVITE message to its SIP server
 - SIP servers tracks the INVITE messages for billing and accounting
 - If one can replay some captured INVITE message, he can make calls at other's expense.



INVITE Replay Billing Attack on AT&T VoIP Subscribers

- SIP authentication does protect the INVITE message from the SIP phone to SIP server
 - Has built-in anti-replay protection
 - Hash (nonce, username, password, realm, SIP-method, request-URI)
 - Correct implementation of existing SIP authentication should prevent INVITE replay
 - Vonage
- Surprisingly, AT&T's callvantage appears vulnerable to INVITE replay
 - We could repeatedly replay the captured, legitimate INVITE (with modification on the SDP part) from our SIP phone one week after the original call
 - All replayed calls have been terminated by the AT&T SIP server after about 3 minutes.

INVITE Replay Billing Attack on AT&T VoIP Subscribers – SIP Message Flow

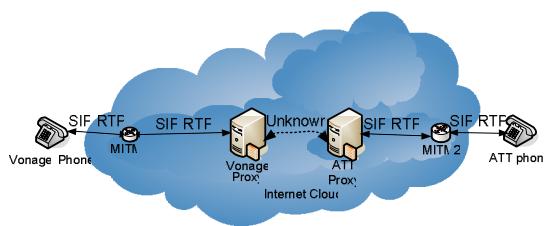


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FakeBusy Billing Attack

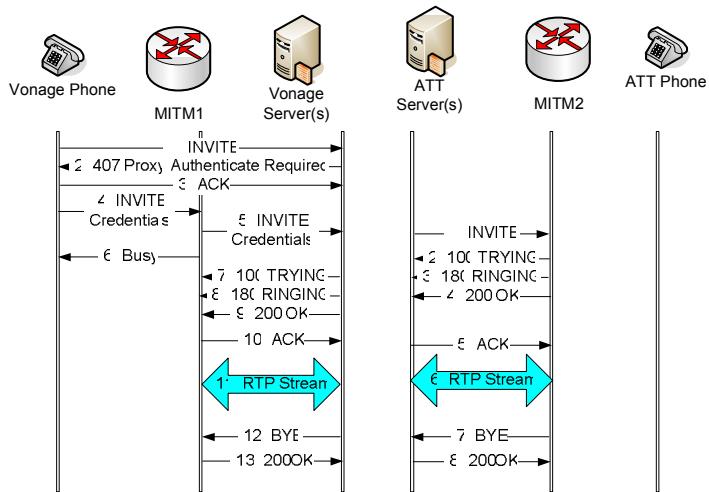
- BUSY message is not protected by SIP authentication
 - The MITM between a SIP phone and SIP server could
 - hijack the VoIP calls of targeted SIP phone
 - Establish the calls with bogus content
 - When both the caller side and callee side have MITM, the MITMs could control the hijacked call duration while keeping the caller and callee unaware of the call
 - This could lead to overcharge on calls the VoIP subscriber has made



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FakeBusy Billing Attack – SIP Message Flow



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ByeDelay and ByeDrop Billing Attack

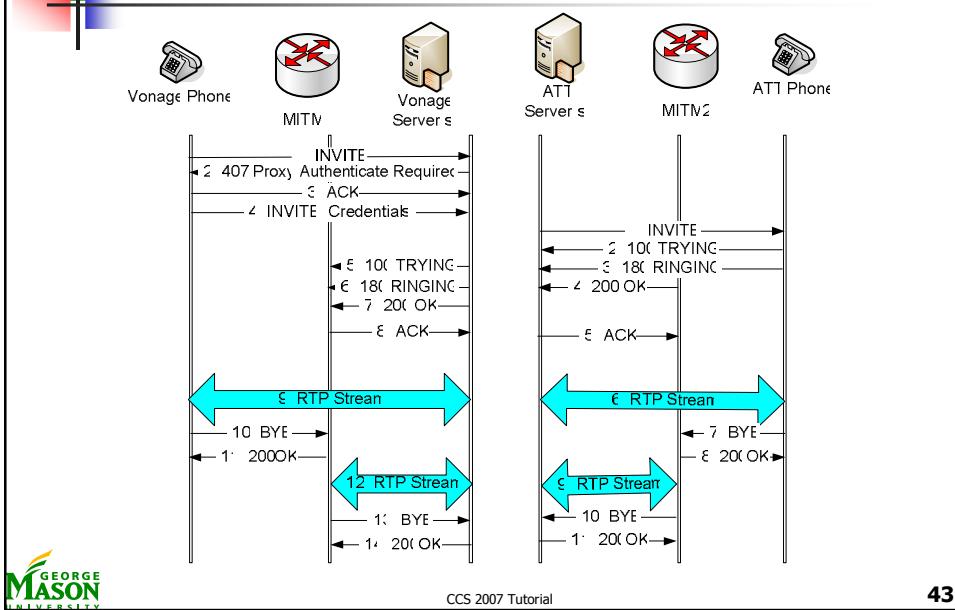
- Established SIP calls are terminated by BYE message
- What if the MITM can delay or simply drop the BYE message of established SIP call?
 - The SIP server will think the call is still alive, and count on the time.
 - This would transparently prolong the duration of established calls
 - This could also lead to overcharge on calls the VoIP subscriber has made
- Such a delay or drop of the BYE message does not involve any modification of the BYE message
 - SIP authentication won't help at all!



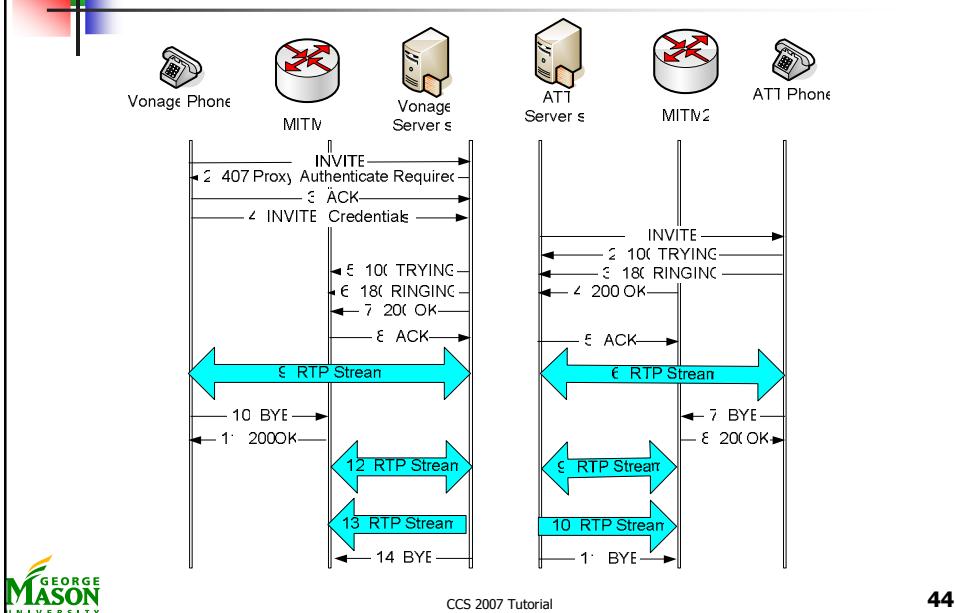
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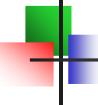
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ByeDelay Billing Attack – SIP Message Flow



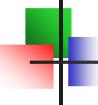
ByeDrop Billing Attack – SIP Message Flow





Potential Mitigations Against VoIP Billing Attacks

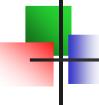
- INVITE replay
 - Just need to implement the SIP authentication correctly
- FakeBusy
 - Simply correlating the SIP and RTP messages won't help
 - Full integrity protection of SIP, RTP could defeat FakeBusy
- ByeDelay, ByeDrop
 - Hard to defend even with full SIP, RTP integrity protection
 - No modification of any SIP or RTP
 - Delay and drop could happen naturally
 - Some heuristics might help



Mitigations against Flooding Type of DoS

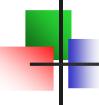
- Measure the difference between the numbers of attempted connections and completed handshakes by Reynolds and Goshal (NDSS 2003)
- Hellinger distance based anomaly detection by Sengar et al (IWQoS 2006)
- SIP-aware firewall by Columbia & Verizon
 - RTP pinhole filtering
 - Return routability check based on null-authentication
 - Could filter out those SIP message with spoofed source
 - State machine sequencing
 - Filter out-of-state SIP messages
 - Maintain dialogue state (source, contact)
 - Only accept BYE with legitimate address





VoIP Intrusion Detection

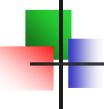
- SCIDIVE by Wu et al. (DSN 2004)
 - Sending fake BYE message to only one endpoint will leave an orphan RTP stream from the other endpoint.
 - Stateful, cross protocol correlation could detect this.
 - What if attacker sends bogus BYE to both endpoints?
 - No orphan RTP stream
 - This could indeed happen naturally if both sides hang up at about the same time!
- Interactive protocol state machine based detection by Sengar et al (DSN 2006)
 - Could detect those attacks that do not follow the SIP state machine
 - What if some attacks do follow the SIP state machine?
 - CallerID spoofing
 - Registration hijacking



Challenges in Securing VoIP

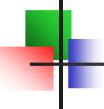
- Open architecture
- Real-time constraints
- Multiple protocols
- Key management issues
- E-911





Challenge #1: Open Architecture

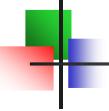
- VoIP network is as open as the Internet
 - The endpoints (SIP phone) of VoIP could be anywhere, and they can freely change their location
- It's likely that we will interact with some endpoint that has NO prior established trust at all
 - We have to allow an unknown person calls us from an unknown SIP phone at an previously unknown IP address
 - We need to be able to call an unknown person at an previously unknown IP address
- How do we authenticate a caller/callee we don't know at all?
- How do we secure VoIP in such a case?
- Trust is NOT transitive!



Challenge #2: Real-Time Constraints

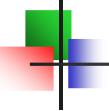
- Unlike other Internet applications such as Web, VoIP has stringent real-time constraints
 - End-to-end delay should be no more than 150ms
 - Each packet of voice stream should be delivered at constant rate (e.g., once every 20ms or 30ms)
 - The processing of each RTP packet should never be more than 20ms
 - The call setup time should not be too long
 - Callers expect to hear ring tone or voice within seconds after dialing
- All these put an upper bound on the total time that all the security mechanisms in all the SIP proxies, RTP servers and SIP phones can use





Challenge #3: Multiple Protocols

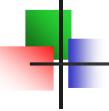
- VoIP involves a number of different protocols
 - Signaling
 - SIP
 - MGCP
 - H.323
 - Voice stream
 - RTP
 - SRTP
 - Security
 - HTTP digest – incorporated in SIP
 - TLS
 - IPsec
 - Key management
- In security, the whole system is as strong as its weakest point!



Challenge #4: Key Management

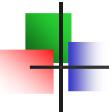
- How could we establish session key with someone we don't know at all?
- Can we expect every SIP phone to have its own private/public keys?
 - How do we authenticate the public key of someone we don't know?
 - Can we expect PKI infrastructure for SIP phones?
- Scalability issue for the SIP servers
 - A SIP server may need to serve tens (or even hundreds) of thousands of subscribers
 - Dynamically establishing and managing keys with many concurrent callers/callees may overwhelm the SIP server





Challenge #5: E-911

- How to provide authenticated caller ID?
 - It's easy to spoof caller ID
- How to determine the real origin of the call?
 - VoIP phones are highly mobile, they can be used from anywhere on the Internet
- How to route a VoIP 911 call to the appropriate PSAP (Public Safety Answering Point)?
- How to keep VoIP 911 connection with PSAP?
 - PSTN 911 call can only be terminated by PSAP
 - VoIP 911 call can be terminated by the caller through power cycle the SIP phone



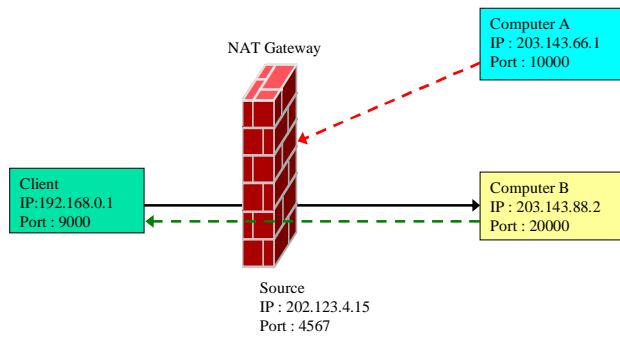
Challenges #6 Firewall Traversal

- Firewall problem with VoIP
 - Firewall needs to block traffic based on
 - Source
 - Destination
 - Traffic type
 - VoIP needs to allow unsolicited incoming calls from unknown and untrusted sources
 - Conflicting requirements
- SIP-aware firewall
 - Allow SIP signaling message and corresponding RTP comes in
- Any exploits of SIP or RTP could compromise the security of SIP-aware firewall
- How to support such VoIP calls without compromising the firewall filtering policy?



Challenges #7 NAT Traversal

- NAT (Network Address Translation)
 - Dynamically maps internal, private IP & port with external, public IP & port
 - Allows multiple hosts in a private network to share one public IP address
 - Most residential IP phones are behind NAT



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Challenges #7 NAT Traversal (cont'd)

- NAT problem with VoIP
 - NAT blocks unsolicited incoming traffic – it does not know how to translate that into private IP & port
 - VoIP needs to support unsolicited incoming calls from previously unknown and untrusted sources
 - SIP phones behind NAT will use its private IP for REGISTER, INVITE and 200OK
 - SIP phones behind NAT will dynamically choose the port number for receiving the incoming RTP stream and specify it in SDP part of INVITE or 200 OK messages



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Challenges #7 NAT Traversal (cont'd)

- Existing NAT traversal solutions
 - UPnP (Universal Plug and Play)
 - Queries the NAT device
 - Does not work with cascading NAT
 - STUN (Simple Traversal of UDP through NAT)
 - Uses TURN server on the public Internet
 - Works with Full Cone, Restricted Cone, and Port Restricted Cone NAT
 - Does NOT work with Symmetric (bidirectional) NAT
 - Does not support TCP
 - SIP RFC 3261 mandates TCP support
 - Susceptible to port scan
 - TURN (Traversal Using Relay NAT)
 - Relies on a TURN server in the middle of the signaling and media path
 - Works with Symmetric (bidirectional) NAT
 - May open door for MITM attack
 - ICE (Interactive Connectivity Establishment)
 - Make use of STUN and TURN



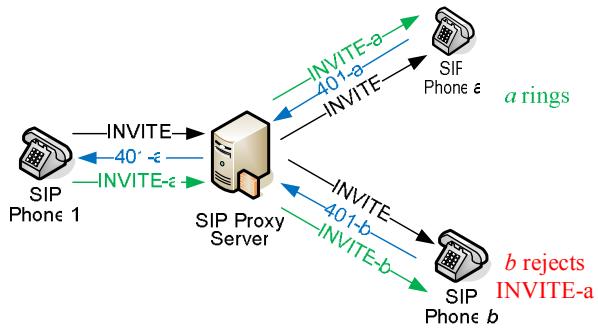
Challenges #7 NAT Traversal (cont'd)

- Security implications of NAT in VoIP
 - Makes it hard to enforce the integrity from end-to-end
 - Makes it hard to validate/authenticate the IP address of SIP phone
 - Opens door for
 - Registration hijacking
 - Call hijacking
 - MITM attack



Challenges #8 Challenge/Response with SIP Forking

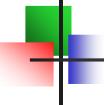
- Challenge/Response would prevent replay attack, but it is difficult to work with SIP forking



Challenges #9 SPIT (Spam on Internet Telephony)

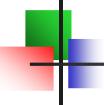
- VoIP spam could be more disturbing than email spam
 - People are used to respond to a phone call immediately
 - What if some spammer keeps ringing your phone all day long?
 - Could make your phone virtually unusable
- VoIP spam is harder to block than email spam
 - Caller ID is not trustworthy
 - Voice content based filtering is difficult





Challenges #9 SPIT (Spam on Internet Telephony) cont'd

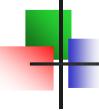
- Turing test against SPIT
 - Before accepts a call, the callee's SIP phone (or proxy) asks the caller some question that is easy for human to answer but difficult for machine.
- However, such a voice Turing test scheme could be abused to launch DoS attack – similar to Smurf attack
 - Attacker send INVITE to lots of VoIP subscribers with spoofed IP address of the victim
 - The proxies of those VoIP subscribers will send RTP streams to the victims
 - The victim could be overwhelmed by those RTP streams



Challenges #10 Voice Phishing (aka Vishing)

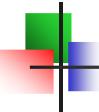
- Attackers started using VoIP in phishing
 - Voice phishing attack on PayPal
<http://www.eweek.com/article2/0,1759,1985966,00.asp>
 - Voice phishing attack on Santa Barbara Bank & Trust
<http://www.websense.com/securitylabs/alerts/alert.php?AlertID=534>
- This essentially exploits people's trust in land line telephone.
- **How to make VoIP as trustworthy as land line telephone?**





References

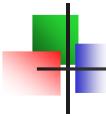
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Summary

- VoIP is complicated, and it involves multiple protocols
- It is very challenging to secure VoIP while keeping it publicly accessible from unknown and untrusted sources
- We want bring the attention of research community to the problems of securing VoIP
- By no means, this tutorial is complete. I hope it will motivate more researchers to look into the security problems of VoIP





Thanks

Questions for me?

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